## Semidefinite hierarchies for polynomial optimization

## CWI



Monique Laurent
Journées SMAI-MODE 2016, Toulouse

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Based on joint works with Etienne de Klerk (Tilburg), Zhao Sun (Montreal), Jean Lasserre, Roxana Hess (Toulouse), Pablo Parrilo (MIT)

## Polynomial optimization



Minimize a polynomial function $f$ over a region

$$
K=\left\{x \in \mathbb{R}^{n}: g_{1}(x) \geq 0, \ldots, g_{m}(x) \geq 0\right\}
$$

defined by polynomial inequalities

$$
\text { Compute: } f_{\min }=\min _{x \in K} f(x)
$$

This is a hard problem, even for simple sets $K$ like

- the standard simplex

$$
\Delta_{n}=\left\{x \in \mathbb{R}^{n}: x_{1}, \ldots, x_{n} \geq 0, \sum_{l=1}^{n} x_{i}=1\right\}
$$

- the hypercube $Q_{n}=[0,1]^{n}$
- the unit sphere

$$
S^{n-1}=\left\{x \in \mathbb{R}^{n}: \sum_{i=1}^{n} x_{i}^{2}=1\right\}
$$

It captures hard combinatorial optimization problems like computing the stability number $\alpha(G)$ and Max-Cut.

## Polynomial optimization formulations for $\alpha(G)$

- Optimization over the simplex:
[Motzkin-Straus 1965]

$$
\frac{1}{\alpha(G)}=\min x^{T}\left(I+A_{G}\right) x \text { s.t. } \sum_{v \in V} x_{v}=1, x_{v} \geq 0(v \in V)
$$

- Optimization over the hypercube:
[Park-Hong 2011]

$$
\alpha(G)=\max \sum_{u \in V} x_{u}-\sum_{u v \in E} x_{u} x_{v} \text { s.t. } x \in[0,1]^{n}
$$

- Optimization over the unit sphere:
[Nesterov 2003]

$$
\frac{2 \sqrt{2}}{3 \sqrt{3}} \sqrt{1-\frac{1}{\alpha(G)}}=\max 2 \sum_{i j \in \bar{G}} z_{i j} y_{i} y_{j} \text { s.t. }(y, z) \in S^{n+m-1}
$$

Lower bounds for polynomial optimization

To approximate:

$$
f_{\min }=\min _{x \in K} f(x), \quad \text { where } K=\left\{x: g_{1}(x) \geq 0, \ldots, g_{m}(x) \geq 0\right\}
$$

## use LP/SDP hierarchies

[Shor (1987), Nesterov (2000), Parrilo, Lasserre (2000-)]

- Express $f_{\min }=\sup \lambda$ s.t. $f(x)-\lambda \geq 0$ over $K$
- Replace nonnegativity by easier sufficient conditions


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Testing whether $f$ is nonnegative is hard, but
testing whether $f$ is a sum-of-squares (SoS): $f=\sum_{j} g_{j}^{2}$
can be done efficiently using semidefinite programming (SDP)

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- Get lower bounds $\underline{f}_{\text {lp }}^{(r)}, \underline{-}_{\text {soss }}^{(r)}, f_{\text {sosP }}^{(r)}$ for $f_{\min }$ by bounding degrees: $\operatorname{deg} g^{\alpha}, \operatorname{deg}\left(\sigma_{J} g^{J}\right), \operatorname{deg}\left(\sigma_{j} g_{j}\right) \leq r$.


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$$
\left\{\underline{f}_{\mathrm{lp}}^{(r)}, \underline{f}_{\mathrm{sosP}}^{(r)}\right\} \leq \underline{f}_{\mathrm{sos} S}^{(r)} \leq f_{\min }
$$

## Representation results for positive polynomials

Theorem
Assume $K$ is compact and $f$ is strictly positive on $K$.

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(1) [Krivine-Handelman] Assume $K$ full-dimensional polytope (all $g_{j}$ have degree 1). Then $f=\sum_{\alpha \in \mathbb{N}^{m}} \lambda_{\alpha} \prod_{j=1}^{m} g_{j}^{\alpha_{j}}$, where $\lambda_{\alpha} \geq 0$.

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- Finite convergence holds generically for SoS bounds.
[Nie 14]
- What about the rate of convergence?


## Rate of convergence of SoS lower bounds

Theorem
Assume $K \subseteq(-1,1)^{n}$. For $f=\sum_{\alpha} f_{\alpha} x^{\alpha}$, set $L_{f}=\max _{\alpha}\left|f_{\alpha}\right| \frac{\alpha \mid}{|\alpha| \mid}$.

- [Schweighofer 2004] Analysis of Schmüdgen type bounds: There exists a constant $c>0$ such that for any polynomial $f$ of degree d:

$$
f_{\min }-\underline{f}_{\text {sosS }}^{(r)} \leq c d^{4} n^{2 d} L_{f} \frac{1}{\sqrt[c]{r}} \quad \text { for } r \geq c d^{c} n^{c d}
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Better results for some simple sets $K$ ?

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Better results for some simple sets $K$ ?
Can choose $c=1$ for Schmüdgen type bounds for simplex \& cube.

Upper bounds for polynomial optimization

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- For simple sets $K=\Delta_{n}$ or $Q_{n}$ : minimize $f$ over the rational grid points with given denominator $r$.


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- For general compact sets $K$, use Lasserre idea:

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Theorem (Lasserre 2011)
For $K$ compact one may use sum-of-squares density functions:

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Bounding degree: $\operatorname{deg}(h) \leq 2 r$, get upper bounds $\bar{f}_{\text {sos }}^{(r)}$ converging to $f_{\text {min }}$.
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## This talk

Analysis of the rate of convergence:

- For the simplex $\Delta_{n}$ and the cube $Q_{n}$ : regular grid upper bounds and the LP lower bounds.
- For $K$ compact: SoS-density upper bounds.
- For the cube $Q_{n}$ : other upper bounds (using other density functions).


# Polynomial optimization OVER THE SIMPLEX 

Upper bounds: For $r \geq 1$

$$
\begin{gathered}
\Delta_{n}(r)=\left\{x \in \Delta_{n}: r x \in \mathbb{N}^{n}\right\} \\
f_{\min , \Delta_{n}(r)}=\min _{x \in \Delta_{n}(r)} f(x)
\end{gathered}
$$

Example: $n=3$ :


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Example: $n=3$ :


$$
r=3
$$

Let $f=\sum_{|\beta|=d} f_{\beta} x^{\beta}$ homogeneous of degree $d$.

## Upper bounds:

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f_{\min } \leq f_{\min , \Delta_{n}(r)}=\min _{x \in \Delta_{n}(r)} f(x)
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Let $f=\sum_{|\beta|=d} f_{\beta} x^{\beta}$ homogeneous of degree $d$.

## Upper bounds:

$$
f_{\min } \leq f_{\min , \Delta_{n}(r)}=\min _{x \in \Delta_{n}(r)} f(x)=\min _{\alpha \in \mathbb{N}^{n}:|\alpha|=r} f(\alpha / r)
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Let $f=\sum_{|\beta|=d} f_{\beta} x^{\beta}$ homogeneous of degree $d$.

## Upper bounds:

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f_{\min } \leq f_{\min , \Delta_{n}(r)}=\min _{x \in \Delta_{n}(r)} f(x)=\min _{\alpha \in \mathbb{N}^{n}:|\alpha|=r} f(\alpha / r) \quad\left(=\sum_{|\beta|=d} f_{\beta} \frac{\alpha^{\beta}}{r^{d}}\right)
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## Lower bounds:

$$
f_{\min } \geq f_{-\mathrm{lp}}^{(r)}=\sup \lambda \text { s.t. } f(x)-\lambda=\underbrace{h(x)}_{\mathbb{R}_{+}[x] r}+\underbrace{u(x)\left(1-\sum_{i=1}^{n} x_{i}\right)}_{\mathbb{R}[x] r}
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\text { where } r^{\underline{d}}=r(r-1) \cdots(r-d+1)
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\underline{f}_{l \mathrm{p}}^{(r)} \leq f_{\min } \leq f_{\min , \Delta_{n}(r)}
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## Error analysis: $\underline{f}_{\mathrm{lp}}^{(r)} \leq f_{\min } \leq f_{\min , \Delta_{n}(r)}$

Theorem (De Klerk-L-Parrilo 2006)

- For any polynomial $f$ of degree $d$ :

$$
f_{\min , \Delta_{n}(r)}-f_{\min }, f_{\min }-\underline{f}_{\mathrm{lp}}^{(r)} \leq \frac{C_{d}}{r}\left(f_{\max }-f_{\min }\right)
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## Error analysis: $\underline{f}_{1 \mathrm{p}}^{(r)} \leq f_{\min } \leq f_{\min , \Delta_{n}(r)}$

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$$
f_{\min , \Delta_{n}(r)}-f_{\min }, f_{\min }-\underline{f}_{\mathrm{lp}}^{(r)} \leq \frac{c_{d}}{r}\left(f_{\max }-f_{\min }\right)
$$

- Can compute the bounds via $\left|\Delta_{n}(r)\right|=O\left(n^{r}\right)$ function evaluations.


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- [De Klerk-L-Sun 2015] May choose $C_{f}=m C_{d}\left(f_{\max }-f_{\min }\right)$, if $f$ has a rational minimizer with denominator $m$.

Key idea for the $1 / r$ convergence rate of $f_{\min , \Delta_{n}(r)}$

- Use the Bernstein approximation of $f$ of order $r$ :

$$
B_{r}(f)(x)=\sum_{\alpha \in \mathbb{N}^{n}:|\alpha|=r} f\left(\frac{\alpha}{r}\right) \frac{r!}{\alpha!} x^{\alpha} .
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$$

- Using properties of Bernstein approximations, one can show:

$$
\max _{x \in \Delta_{n}} B_{r}(f)(x)-f(x) \leq \frac{C_{d}}{r} .
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[De Klerk-L-Sun 2014]

Polynomial optimization OVER THE HYPERCUBE

## Upper bounds:

$$
f_{\min } \leq f_{\min , Q_{n}(r)}=\min _{x \in Q_{n}(r)} f(x)=\min _{\alpha \in \mathbb{N}^{n}: \alpha_{i} \leq r} f(\alpha / r)
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Measure-Based upper Bounds

## Upper bounds using SoS densities

Lasserre (2011) proved: For K compact

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\begin{aligned}
& f_{\min }=\min _{x \in K} f(x)=\min _{\mu \text { probability measure on } K} \int_{K} f(x) d \mu \\
& =\inf \int_{K} f(x) h(x) d x \text { s.t. } h \operatorname{SoS}, \int_{K} h(x) d x=1
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## Theorem (De Klerk-L-Sun 2015)

Assume $K$ is compact and 'nice' (e.g. convex body) and $f$ has Lipschitz constant $M_{f}$. There exist constants $C_{K}>0$ and $r_{K} \geq 1$ such that

$$
\bar{f}_{\text {sos }}^{(r)}-f_{\min } \leq \frac{C_{K} M_{f}}{\sqrt{r}} \quad \forall r \geq r_{K} .
$$

Example: Motzkin polynomial on $K=[-2,2]^{2}$

$$
f\left(x_{1}, x_{2}\right)=x_{1}^{4} x_{2}^{2}+x_{1}^{2} x_{2}^{4}-3 x_{1}^{2} x_{2}^{2}+1
$$

Global minimizers: $(-1,-1),(-1,1),(1,-1),(1,1)$.


## Example: Motzkin polynomial on $[-2,2]^{2}$ (ctd.)

Optimal SoS density $h$ of degree 12


## Example: Motzkin polynomial on $[-2,2]^{2}$ (ctd.)

Optimal SoS density hof degree 16


## Example: Motzkin polynomial on $[-2,2]^{2}$ (ctd.)

Optimal SoS density hof degree 20


## Example: Motzkin polynomial on $[-2,2]^{2}$ (ctd.)

Optimal SoS density hof degree 24


## Convergence analysis: sketch of proof

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Main result: If $f$ has Lipschitz constant $M_{f}$ and $K$ is nice at a then

$$
\int_{K} f(x) c_{\mathbf{a}}^{r} H_{r, \mathbf{a}}(x) d x-f_{\min } \leq \frac{C_{K} M_{f}}{\sqrt{r}}
$$

OTHER MEASURE-BASED UPPER BOUNDS FOR THE HYPERCUBE:

- Handelman type Densities
- SChMÜDGEN TYPE DENSITIES


## Using Handelman type densities for $K=[0,1]^{n}$

For $K=[0,1]^{n}$, consider the upper bound:

$$
\bar{f}_{\mathrm{lp}}^{(r)}=\min \int_{K} f(x) h(x) d x \text { s.t. } h(x)=\sum_{\alpha, \beta \in \mathbb{N}^{n}} \underbrace{\lambda_{\alpha, \beta}^{\geq 0} x^{\alpha}(1-x)^{\beta}}_{\text {degree }=r}, \int_{K} h(x) d x=1
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Theorem (De Klerk-Lasserre-L-Sun 2015)

- $\bar{f}_{\mathrm{lp}}^{(r)}$ needs $O\left(n^{r}\right)$ elementary computations:

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\bar{f}_{\mathrm{lp}}^{(r)}=\min _{|\alpha+\beta|=r} \frac{\int_{K} f(x) x^{\alpha}(1-x)^{\beta} d x}{\int_{K} x^{\alpha}(1-x)^{\beta} d x}
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\bar{f}_{\mathrm{lp}}^{(r)}-f_{\min } \leq \frac{C_{f}}{\sqrt{r}}, \quad \leq \frac{C_{f}^{\prime}}{r} \quad \text { if } f \text { has a rational minimizer. }
$$

$\rightsquigarrow$ Link to the beta distribution

Motzkin polynomial:
$f(x, y)=(4 x-2)^{4}(4 y-2)^{2}+(4 x-2)^{2}(4 y-2)^{4}-3(4 x-2)^{2}(4 y-2)^{2}+1$ over $K=[0,1]^{2}$ : Handelman-type densities (deg 24,50) \& SOS (deg 24)




## Using Schmüdgen type densities for $K=[-1,1]^{n}$

For $K=[-1,1]^{n}$ consider the upper bound:

$$
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- Convergence rate: For $r \geq r_{0}$

$$
\bar{f}_{\text {sosS }}^{(r)}-f_{\min } \leq \frac{C_{f}}{r^{2}} .
$$

Motzkin polynomial: $f(x, y)=64\left(x^{4} y^{2}+x^{2} y^{4}\right)-48 x^{2} y^{2}+1$ over $K=[-1,1]^{2}$

Optimal Schmüdgen type densities of degree 12, 16:



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$\rightsquigarrow$ rate of convergence in $1 / r^{2}$ for $\bar{f}_{\text {sosS }}^{(r)}$.

Kernel approximations $\delta_{0}^{(r)}(x)$ of the Dirac at degree $r=8,16,32,64$ :


$$
\frac{\delta_{0}^{(r)}(x)}{\pi \sqrt{1-x^{2}}} \sim \frac{1}{\sqrt{2 \pi \sigma^{2}}} \exp \left(-\frac{x^{2}}{2 \sigma^{2}}\right), \quad \text { with } \sigma \sim \frac{\pi}{r+2}
$$

## Concluding remarks

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| $r$ | Matyas |  | Three-Hump | Camel | Motzkin |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Value | Time (sec.) | Value | Time (sec.) | Value | Time (sec.) |
| 2 | 8.26667 | 0.000739 | 265.774 | 0.000742 | 4.2 | 0.000719 |
| 6 | 4.28172 | 0.000072 | 29.0005 | 0.000066 | 1.06147 | 0.000080 |
| 12 | 2.99563 | 0.000263 | 4.43983 | 0.000263 | 0.801069 | 0.000208 |
| 18 | 1.83356 | 0.000655 | 2.55032 | 0.000586 | 0.565553 | 0.000766 |
| 24 | 1.11785 | 0.001753 | 1.2775 | 0.001693 | 0.406076 | 0.001712 |
| 30 | 0.8524 | 0.002270 | 1.0185 | 0.002936 | 0.3004 | 0.002351 |
| 36 | 0.5760 | 0.005510 | 0.7113 | 0.004882 | 0.2300 | 0.006060 |
| 40 | 0.4815 | 0.006975 | 0.6064 | 0.007031 | 0.1817 | 0.007686 |

Matyas: $f=\left(x_{1}+2 x_{2}-7\right)^{2}+\left(2 x_{1}+x_{2}-5\right)^{2}, \quad K=[-10,10]^{2}$.
Three-Hump Camel: $f=2 x_{1}^{2}-1.05 x_{1}^{4}+\frac{1}{6} x_{1}^{6}+x_{1} x_{2}+x_{2}^{2}, \quad K=[-5,5]^{2}$.
Motzkin: $f=x_{1}^{4} x_{2}^{2}+x_{1}^{2} x_{2}^{4}-3 x_{1}^{2} x_{2}^{2}+1, \quad K=[-2,2]^{2}$.
$f_{\min }=0$, bounds $\bar{f}_{\text {sos }}^{(r)}$ with SoS density.

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Better theoretical convergence results for the upper bounds, but slower in practice...

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Better theoretical convergence results for the upper bounds, but slower in practice...

- Problem: Show better convergence rates for the lower bounds.

New techniques needed ...

## Concluding remarks

- One can use the upper bounds to generate good feasible solutions using sampling.

| $r$ | $\bar{f}_{\text {sos }}^{(r)}$ | Mean | Variance | Minimum | Sample Size |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | 265.774 | 216.773 | 177142.0 | 0.106854 | 20 |
|  |  | 261.23 | 193466.0 | 0.11705 | 1000 |
| 4 | 29.0005 | 28.0344 | 2964.85 | 1.1718 | 20 |
|  |  | 6712.8 | 0.014255 | 1000 |  |
| 14 | 4.43983 | 3.96711 | 20.3193 | 0.260331 | 20 |
|  |  | 3.7911 | 57.847 | 0.0076111 | 1000 |
| 22 | 1.71275 | 1.30757 | 1.90985 | 0.0320489 | 20 |
|  |  | 1.6379 | 7.2518 | 0.0021144 | 1000 |
| 24 | 1.27749 | 0.841194 | 0.914514 | 0.0369565 | 20 |
|  |  | 2.3 | 0.0005154 | 1000 |  |
| Uniform Sample | 304.032 | 163021.0 | 1.65885 | 20 |  |
|  | 243.216 | 183724.0 | 0.00975034 | 1000 |  |

SoS upper bounds for the Three-Hamp Camel function:
$f=2 x_{1}^{2}-1.05 x_{1}^{4}+\frac{1}{6} x_{1}^{6}+x_{1} x_{2}+x_{2}^{2}$ over $K=[-5,5]^{2}$.

## Thank you

## Based on the papers

- Improved convergence rates for Lasserre-type hierarchies of upper bounds for box-constrained polynomial optimization. With E. de Klerk, R. Hess, arXiv:1603.03329.
- Bound-constrained polynomial optimization using only elementary calculations. With E. de Klerk, J. Lasserre and Z. Sun, arXiv:1507.04404.
- Convergence analysis for Lasserre's measure-based hierarchy of upper bounds for polynomial optimization. With E. de Klerk and Z. Sun, arXiv:1411.6867
- On the convergence rate of grid search for polynomial optimization over the simplex. With E. de Klerk, Z. Sun, J. Vera, Opt. Letters, 2016.
- An alternative proof of a PTAS for fixed-degree polynomial optimization over the simplex. With E. de Klerk and Z. Sun. Math.Prog. 2014.
- A PTAS for the minimization of polynomials of fixed degree over the simplex. With E. De Klerk and P. Parrilo. TCS, 2006.


## Optimization over the unit sphere

SOS lower bounds:

$$
\underline{f}_{\text {sosS }}^{(r)}=\sup \lambda \text { s.t. } f(x)-\lambda=\underbrace{\sigma_{0}}_{\text {SOS of degree } 2 r}+u\left(1-\sum_{i=1}^{n} x_{i}^{2}\right) .
$$

Theorem
Let $f$ be a homogeneous polynomial of even degree.

- [Faybusovich 2003] $f_{\min }-f_{\text {soss }}^{(r)}=O\left(\frac{1}{r}\right)$ for $r=O(n)$.
- In addition, Doherty-Wehner (2013) construct measure-based upper bounds with the same performance guarantee.
- Parrilo-Wehner announced convergence in $O\left(\frac{1}{r^{2}}\right)$ for both upper and lower bounds.

